

Teaching the stack for fun and profit

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A HackerWeek presentation

Me

- TCD
 - Grand
- DCU
 - Grand

Warnings

- Lecture pertains to IA32 for the most part
- No insta-I33t
 - Frames, pointers, bounds
 - Stack discipline is booring
 - Complex topic

Apology before one is required

- Target audience

The stack

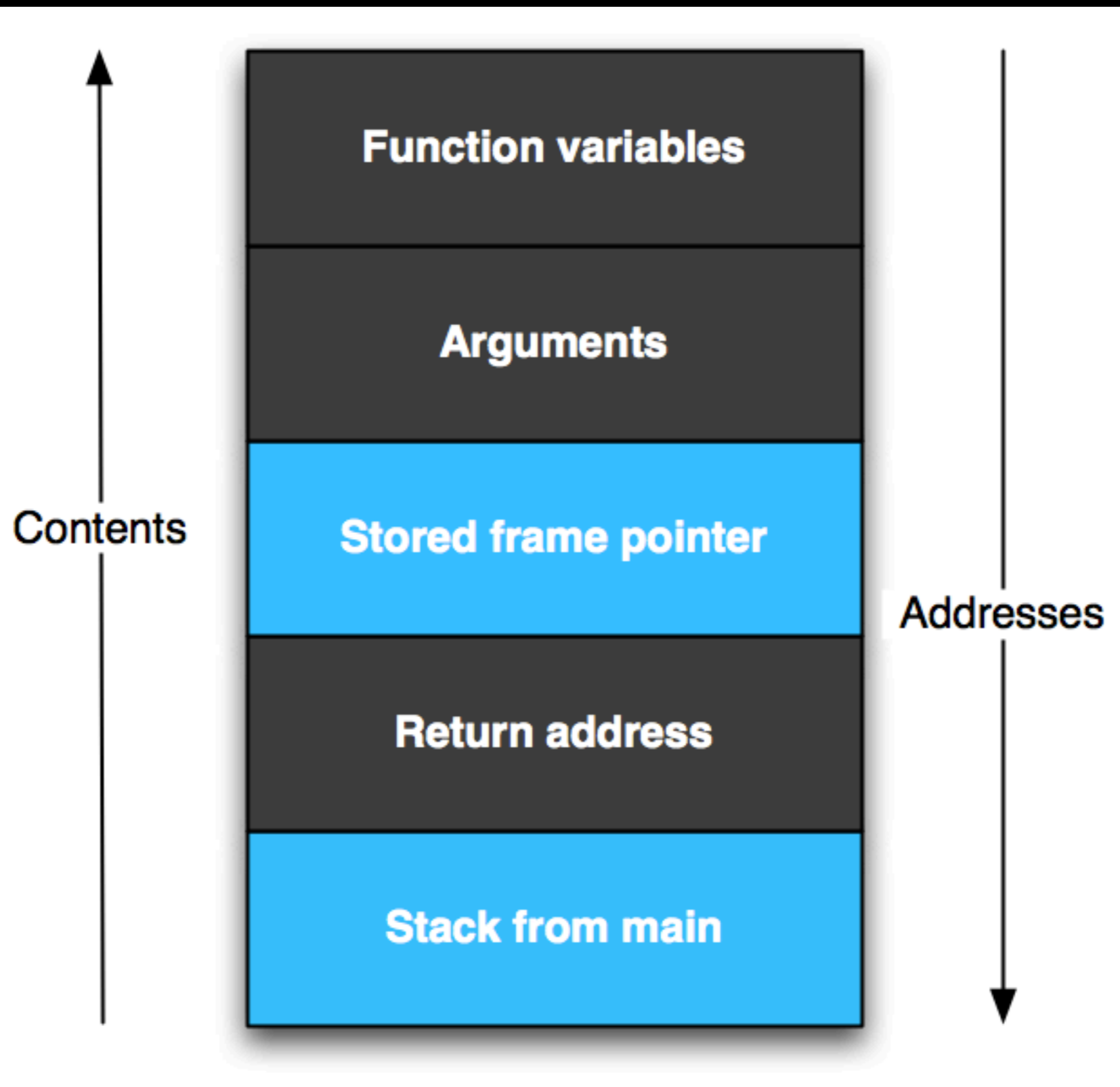
- Keep track of functions called
 - Arguments
 - *Where to return to*
 - Local variables
- Stored in frames

Stack buffer overflows

- Debilitating
- Widespread (still)
- Easily introduced
- Successful exploitation alters control flow
- Flow altered via modification of return values

Function execution

- Before function execution
 - Address of next instruction stored
 - Frame pointer stored
 - Frame pointer becomes function address
- After function execution - ret
 - Return address popped from stack



Bounds

- User controlled input = potential threat
 - SQL injection, XSS, BoFs
- Programming practice should protect
 - Some core elements are insecure
- Input should go into input buffers
 - Nowhere else...

C checklist

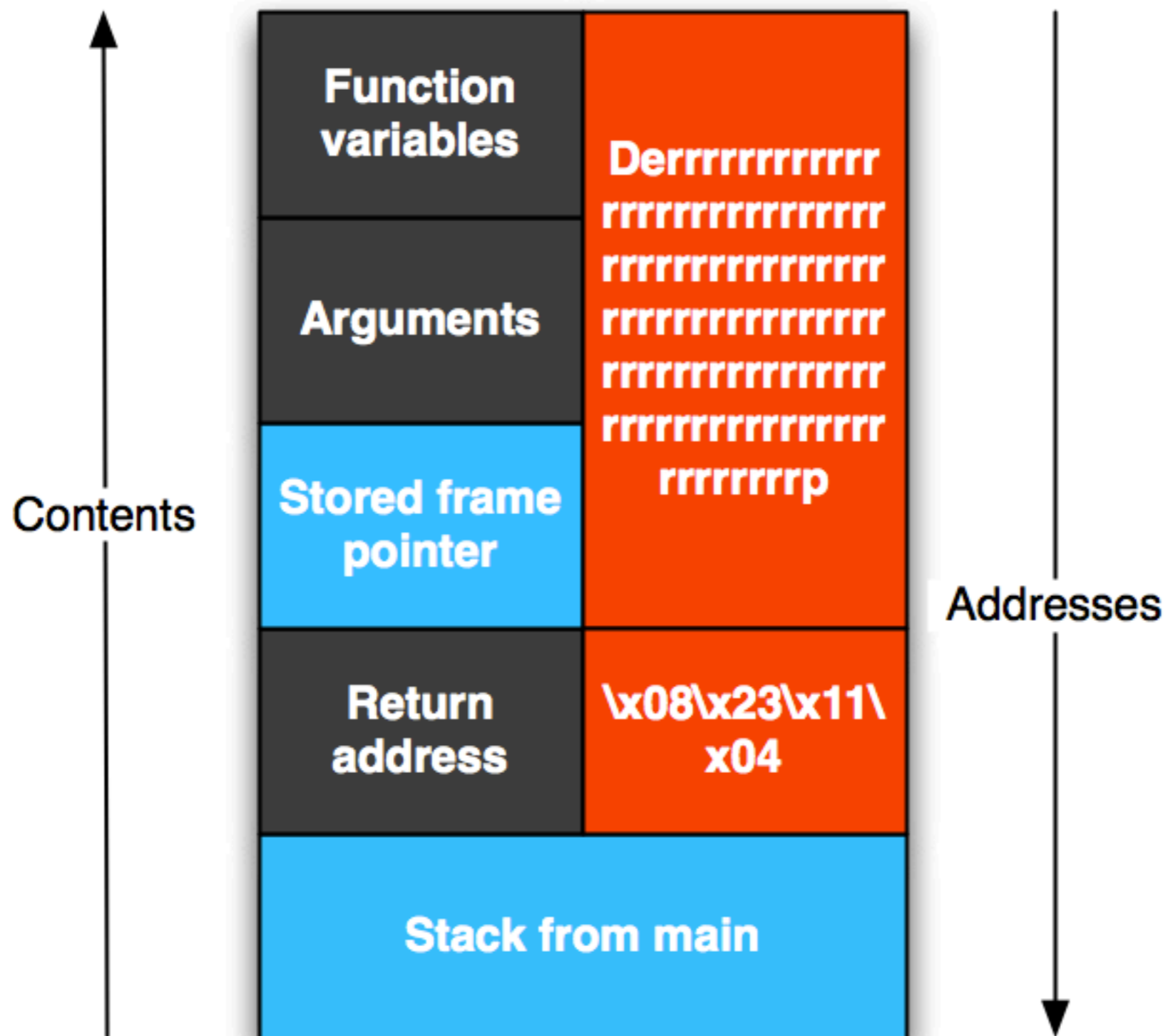
- `gets`
- `printf`
- `strcpy`
- `memcpy`
- Others...

Unchecked bounds

- Unchecked bounds can mean compromise
- What happens here? Inside function `x`:
 - `input` = pointer to 18 bytes of data
 - `destination` = 10 byte array local to `x`
 - `strcpy(destination, input)`

What now?

- Unsurprisingly, a crash
 - Function postlude goes ahead unaware
 - Stored frame pointer becomes “rrrp”
 - Derp
- Let's boost it a little



Exploitation at last

- Input is simply a `\x` escaped hex memory address
- This points to a function from which to return to
- Handy if you have the right function
 - No arguments, calls `execl("/bin/sh",null);`

Best-case

- Assuming this function
- Get the address (say 0x04112308)
- Exploited function executes “ret”
- Code “returns” to the exploit
- Usually not an option

Shellcode

- Expansion of previous technique
- Injection of code alongside memory redirection
- Encoded assembly instructions to execute exploit code
- The shorter your code, the leeter your skilz

Pre-assembly

```
void shellcode(void) {  
    char *execarg[2];  
    execarg[0] = "/bin/sh";  
    execarg[1] = NULL;  
    execve(execarg[0], execarg, NULL);  
}
```

Preparing payloads

- Injecting the code isn't enough
- NOPs used to make up loading space
- Buffer location is unclear to attacker

Assembly

- Produced assembly is optimised
- `\xeb\x1f\x5e\xb8\x00\x00\x00\x00\xc7`
`\x46\x07\x00\x00\x00\x00\x50...`
- Nulls need to be removed
- Strings need to be loaded carefully

Caveats

- Executing `/bin/sh` runs as exploiting user
- SUID bit changes this
 - Program executes as owner
 - Usually aids less-privileged users
 - `chmod +s aprogram`

Prevention

- Writing safe code
- Non-executable stack
- Bounds checking
- Stack Guard
- Instruction randomisation

Correct code

- Easier than it sounds
- Cliches appear time and time again
- Static analysis can help

Bounds checking

- Ensure reads and writes are in-bounds
- Only array references
- Can't see variables beyond where they are declared
- Some operations slowed by large factors

Randomised instructions

- Encrypted executables
- On the fly decryption
- Most injected code lacks valid instructions
- Only protects against injected code
- Performance hit

Canaries

- Present in GCC 4.x
- Augmented function prologues and epilogues
- Canary near return address
- Value integrity is checked on return

No-exec stack

- Non-executable segments in stack
- Disable execute on all but `.text`
 - Code section
- Still vulnerable to
 - Data modification
 - Existing dangerous code

Prognosis

- Grim
- Only you etc. etc.

Further reading

- “Smashing the Stack for Fun and Profit”
 - Aleph One
- Real world examples!
 - exploit-db
 - Shellcode archive

Reading is only so entertaining

- intruded.net
 - Great for learning
- overthewire.org
 - A little more advanced

kthxbai

- Questions?
- Presentation available soon
- <http://www.netsoc.tcd.ie/~nosmo>